

Ultrafilters, congruences, and profinite groups

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based on joint work with
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- In this talk, we mainly adopt the latter.

Definable types

Review of the basics

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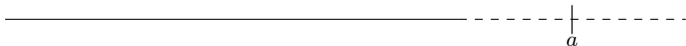
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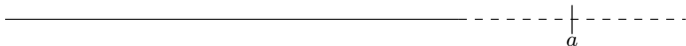
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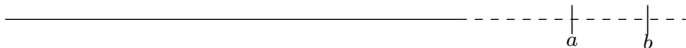
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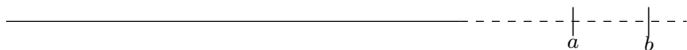
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- $p \otimes q$ is itself definable and \otimes is associative.

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- Corollary: q nonrealised, $(a, b) \models p \otimes q, (b, c) \models q \otimes r \implies (a, c) \models p \otimes r$.

FALSE in general, e.g. random graph

(note: historically \otimes, \oplus were first defined on ultrafilters)

- Spelling out \oplus , with $+$ the usual sum, $A \in p \oplus q \iff \{n : A - n \in q\} \in p$

A typical application in Ramsey theory/additive combinatorics:

Theorem (Hindman)

\forall finite colouring of \mathbb{Z} , there are $(z_i)_{i < \omega}$ s.t. all $z_{i_0} + \dots + z_{i_n}$ have the same colour.

Proof. Use Ellis theory to find $u \in \beta\mathbb{Z}$ with $0 < u = u \oplus u$. Pick the colour $A \in u$. Since $u = u \oplus u$, there is $z_0 \in A$ with $A \cap (A - z_0) \in u$. Again since $u = u \oplus u$, there is $z_1 \in A \cap (A - z_0)$ with $A \cap (A - z_0) \cap (A - z_1) \cap (A - (z_0 + z_1)) \in u$. Repeat. \square

Want finite products instead? Work in $(\beta\mathbb{Z}, \odot)$.

Strong congruence

- Recap: $(\beta\mathbb{Z}, \oplus, \odot)$, applications in Ramsey theory/combinatorics.
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Write $u \equiv_w^s v$ iff for some/all $(d, a, b) \models w \otimes u \otimes v$ we have $d \mid (a - b)$.

(equivalently: $u \equiv_w^s v$ iff $\{n : n\mathbb{Z} \in u \ominus v\} \in w$)

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Maybe there is a better notion of congruence?

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If $\text{BD}(A) > 0$ and $A^{\mathbb{G}}$ is thick, there are $u, v \in \beta\mathbb{Z} \setminus \mathbb{Z}$ such that $A \in u \oplus v$ and $A^{\mathbb{G}} \in v \oplus u$.

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- Failure of transitivity: take any infinite w containing all $n\mathbb{Z} + 1$. Then $0 \equiv_w w \equiv_w 1$.

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Theorem (DLMPR)

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Examples and non-examples

✗ Infinite primes: $w(x) \otimes w(y) \vdash |x| < |y|$ are primes!

Recall w is self-divisible iff, equivalently:

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- ✗ For other *supernatural numbers* (functions $\mathbb{P} \rightarrow \omega + 1$)(formal products $\varphi = \prod_{p \in \mathbb{P}} p^{\varphi p}$, if you prefer)
there are both self-divisible w and non-self-divisible w with that divisibility.

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- ✗ Infinite primes: $w(x) \otimes w(y) \vdash |x| < |y|$ are primes!
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- ✓ Powers of 2! (2 factorial, if you wish)
- ✓ Stuff divisible by every $n > 0$;
e.g. all \oplus -idempotents and all \odot -minimals
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w is self-divisible iff $\ker \sigma_w$ is closed iff $\beta\mathbb{Z} / \equiv_w^s$ is profinite iff $\beta\mathbb{Z} / \equiv_w^s \cong \prod_{p \in \mathbb{P}} G_{p,w}$

$$G_{p,w} = \begin{cases} \mathbb{Z}/p^n\mathbb{Z}, & \text{if } n = \max\{k : p^k\mathbb{Z} \in w\} \\ \mathbb{Z}_p, & \text{otherwise.} \end{cases}$$

The full list

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$$D(w) := \{n : n\mathbb{Z} \in w\}$$

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Theorem 3.10. For every $w \in \beta\mathbb{Z} \setminus \{0\}$, the following are equivalent.

- (1) The ultrafilter w is self-divisible.
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- (5) For every $a, b \models w$ there is $c \models w$ such that $c \mid \gcd(a, b)$.

Theorem 6.8. The following are equivalent for $w \in \beta\mathbb{Z} \setminus \{0\}$.

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- (8) For all v , if $w \equiv_v 0$ then $w \equiv_v^s 0$.
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Thanks for listening!



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