Computational Aspects of Line and Toric Arrangements

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Hyperplane arrangements

Definition (Hyperplane arrangement)

Let V be a finite-dimensional vector space over a field \mathbb{K} . A hyperplane arrangement \mathcal{A} is a (finite) collection of affine hyperplanes of V. The same definition can be given for a projective hyperplane arrangement in a projective space.



Basic definitions

• The complement of an arrangement \mathcal{A} is the set

$$\mathfrak{M}(\mathcal{A})\coloneqq V\smallsetminus \bigcup_{\mathsf{H}\in\mathcal{A}}\mathsf{H}.$$

• An arrangement \mathcal{A} is central if

$$\bigcap_{\mathsf{H}\in\mathcal{A}}\mathsf{H}\neq\varnothing.$$

• The defining polynomial of an arrangement \mathcal{A} is

$$Q_{\mathcal{A}} = \prod_{H \in \mathcal{A}} \alpha_H$$

where α_H is a linear form defining H.

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Intersection poset

Definition (Intersection poset)

The intersection poset L(A) of an arrangement A is the set of all non-empty intersections of hyperplanes of A, partially ordered by reverse inclusion. It includes V as the intersection of zero hyperplanes.



A, B, C, D are the singular points of \mathcal{A} . For a singular point P, its *multiplicity* m(P) is the number of lines passing through it.



Combinatorial properties

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We say that a property of an arrangement A is combinatorial if it depends only on the intersection poset L(A).



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- ► The cohomology ring H*(M(A); C) is combinatorial (Orlik-Solomon algebra).
- ► The fundamental group $\pi_1(\mathcal{M}(\mathcal{A}))$ is not combinatorial (Rybnikov counterexample).

Local systems

Definition (Local system)

Let \mathcal{A} be an arrangement of n hyperplanes in \mathbb{C}^m , $M := \mathcal{M}(\mathcal{A})$, and let R be a commutative ring with unity. A rank-1 local system is a structure of $\pi_1(M)$ -module on R.



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When $R = \mathbb{C}$, the action $\pi_1(M) \to \mathcal{A}ut(\mathbb{C}) \simeq \mathbb{C}^*$ factors through $H_1(M; \mathbb{Z})$, which is free abelian of rank n generated by β_1, \ldots, β_n , where β_i is a loop around a line of \mathcal{A} . In this case, the local system is defined by a choice of a non-zero complex number t_i for each β_i .



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We will denote by \mathbb{C}_t the local system defined by $t \coloneqq (t_1, \ldots, t_n) \in (\mathbb{C}^*)^n$, and with $H_*(M; \mathbb{C}_t)$ and $H^*(M; \mathbb{C}_t)$ respectively the homology and cohomology with coefficients in \mathbb{C}_t .



Characteristic varieties

Definition (Characteristic variety)

Let $\mathcal A$ be an arrangement as before. The (first) characteristic variety is

 $\mathcal{V}(\mathcal{A}) \coloneqq \{ \mathbf{t} \in (\mathbb{C}^*)^n \mid \dim H^1(M; \mathbb{C}_{\mathbf{t}}) \geqslant 1 \}.$



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Theorem (Arapura '97)

 $\mathcal{V}(\mathcal{A})$ is a union of (possibly translated) subtori of $(\mathbb{C}^*)^n$.



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Theorem (Arapura '97)

 $\mathcal{V}(\mathcal{A})$ is a union of (possibly translated) subtori of $(\mathbb{C}^*)^n$.

Is $\mathcal{V}(\mathcal{A})$ combinatorial?



Let A be the Orlik-Solomon algebra associated with A. Fix $a \in A^1$. Left-multiplication by a gives A^{\bullet} the structure of a cochain complex.

Definition (Resonance variety)

The (first) resonance variety is

$$\mathcal{R}(\mathcal{A}) \coloneqq \{ a \in \mathcal{A}^1 \mid \dim H^1((\mathcal{A}^{\bullet}, a \cdot); \mathbb{C}) \ge 1 \}.$$

 $\mathfrak{R}(\mathcal{A})$ is a union of linear subspaces of $A^1\simeq \mathbb{C}^n$



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Definition (Resonance variety)

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 $\mathfrak{R}(\mathcal{A})\coloneqq \{ \mathfrak{a}\in A^1 \mid \dim H^1((A^\bullet,\mathfrak{a}\,\cdot\,);\mathbb{C})\geqslant 1 \}.$

 $\mathfrak{R}(\mathcal{A})$ is a union of linear subspaces of $A^1\simeq \mathbb{C}^n$

Tangent Cone Theorem (Cohen-Suciu '99)

 $\mathfrak{R}(\mathcal{A})$ is the tangent cone of $\mathcal{V}(\mathcal{A})$ at $(1,\ldots,1)\in (\mathbb{C}^*)^n.$

The "homogeneous part" of $\mathcal{V}(\mathcal{A})$ is combinatorial!



From now we will suppose that \mathcal{A} is an arrangement of n + 1 projective lines in $\mathbb{P}^2(\mathbb{R})$. The defining polynomial $Q_{\mathcal{A}}$ belongs to $\mathbb{R}[X, Y, Z]$ and it is homogeneous of degree n + 1.

Since the topology of $\mathcal{M}(\mathcal{A})$ in $\mathbb{P}^2(\mathbb{R})$ is easy to describe, we will consider the *complexified* arrangement $\mathcal{A}_{\mathbb{C}}$, which is the arrangement in $\mathbb{P}^2(\mathbb{C})$ defined by $Q_{\mathcal{A}}$, and study the complement $\mathcal{M}(\mathcal{A}_{\mathbb{C}}) \subseteq \mathbb{C}^2$.



If ${\mathcal A}$ is an arrangement of n+1 projective lines in ${\mathbb P}^2({\mathbb C}),$ it is known that

$$\mathrm{H}_1(\mathrm{\mathfrak{M}}(\mathcal{A});\mathbb{C}) = \left<\beta_1,\ldots,\beta_{n+1} \mid \beta_1\cdots\beta_{n+1} = 1\right>$$

with the commutation relations.



Our setting

If ${\mathcal A}$ is an arrangement of n+1 projective lines in ${\mathbb P}^2({\mathbb C}),$ it is known that

$$H_1(\mathcal{M}(\mathcal{A});\mathbb{C}) = \langle \beta_1, \dots, \beta_{n+1} \mid \beta_1 \cdots \beta_{n+1} = 1 \rangle$$

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1. If (t_1, \ldots, t_{n+1}) is a local system for \mathcal{A} , then $t_1 \cdots t_{n+1} = 1$.



Our setting

If ${\mathcal A}$ is an arrangement of n+1 projective lines in ${\mathbb P}^2({\mathbb C}),$ it is known that

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with the commutation relations.

- 1. If (t_1, \ldots, t_{n+1}) is a local system for \mathcal{A} , then $t_1 \cdots t_{n+1} = 1$.
- 2. Let $a\mathcal{A}$ be the arrangement of n lines in \mathbb{C}^2 obtained by sending ℓ_{n+1} to infinity. Then

$$\mathcal{V}(\mathcal{A}) = \{(\mathbf{t}, \mathbf{t}_{n+1}) \in (\mathbb{C}^*)^{n+1} \mid \mathbf{t} \in \mathcal{V}(a\mathcal{A}), \ \mathbf{t}_1 \cdots \mathbf{t}_{n+1} = 1\}.$$

Local and non-local components

Denote the lines of \mathcal{A} with $[n + 1] \coloneqq \{1, \ldots, n + 1\}$ and a singular point with the subset of [n + 1] indicating the lines passing through it. Let $S \subseteq \mathscr{O}([n + 1])$ be the set of the singular points.

For each $P\in S$ with $\#(P)\geqslant 3,$ there is a local component of $\mathcal{R}(\mathcal{A})$ given by

$$C(\mathbf{P}) \coloneqq \left\{ z \mid \sum_{j=1}^{n+1} z_j = 0 \right\} \cap \bigcap_{j \notin \mathbf{P}} \{ z \mid z_j = 0 \}$$

The *non-local* components admit a description in terms of *neighbourly partitions*.

Neighbourly partitions

Definition (Neighbourly partition)

A partition $\pi=(p_1\mid\cdots\mid p_r)$ of [n+1] is neighbourly if for all $i=1,\ldots,r$ and for all $P\in S$

 $\texttt{\#}(p_i \cap P) \geqslant \texttt{\#}(P) - 1 \Rightarrow P \subseteq p_i.$

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$$\#(p_i \cap P) \geqslant \#(P) - 1 \Rightarrow P \subseteq p_i.$$

If π is a neighbourly partition, define $C(\pi) \subseteq \mathbb{C}^{n+1}$ as

$$C(\pi) \coloneqq \left\{ z \; \left| \; \sum_{j=1}^{n+1} z_j = 0 \right\} \cap \bigcap_{P \in \mathcal{P}} \left\{ z \; \left| \; \sum_{j \in P} z_j = 0 \right\} \right.$$

where $\mathcal{P} \coloneqq \{P \in S \; | \not \exists \; p \in \pi \text{ s.t. } P \subseteq p\}.$

Neighbourly partitions

Proposition

If dim $(C(\pi)) \ge 2$, then $C(\pi)$ is a non-local component of $\mathcal{R}(\mathcal{A})$.

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Proposition

If dim $(C(\pi)) \ge 2$, then $C(\pi)$ is a non-local component of $\mathcal{R}(\mathcal{A})$.

If π is a partition of a subset $B \subseteq [n + 1]$, define support of π , supp (π) , the set B.

Proposition

Let $\mathcal{B} \subseteq \mathcal{A}$ be a subarrangement and let π be a neighbourly partition for \mathcal{B} such that dim $(C(\pi)) \ge 2$. Then

$$C(\pi)\cap \bigcap_{j\notin \mathsf{supp}(\pi)}\{z_j=0\}$$

is a non-local component of $\mathcal{R}(\mathcal{A}).$ All non-local components of $\mathcal{R}(\mathcal{A})$ arise from subarrangements of $\mathcal A$ this way.

Combinatorics of the characteristic variety

For the homogeneous part of the characteristic variety $\mathcal{V}(\mathcal{A})$, we define *ideals* of $\mathbb{C}[\mathsf{T}_1^{\pm 1}, \ldots, \mathsf{T}_{n+1}^{\pm 1}]$ such that their varieties are the components of $\mathcal{V}(\mathcal{A})$.

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• If $P \in S$ with $\#(P) \ge 3$, define

$$\mathbb{J}(\mathsf{P})\coloneqq \left(\prod_{j=1}^{n+1}\mathsf{T}_j-1\right)+\big(\mathsf{T}_j-1\bigm| j\notin\mathsf{P}\big);$$

this corresponds to a local component of $\mathcal{V}(\mathcal{A})$.



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• If $P \in S$ with $\#(P) \ge 3$, define

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this corresponds to a local component of $\mathcal{V}(\mathcal{A}).$

• If π is a neighbourly partition, define

$$\mathbb{I}(\pi)\coloneqq \left(\prod_{j=1}^{n+1}T_j-1\right)+\left(\prod_{j\in P}T_j-1\;\middle|\;P\in\mathcal{P}\right)$$

where $\mathcal{P} := \{ P \in S \mid \nexists \ p \in \pi \text{ s.t. } P \subseteq p \}.$

Combinatorics of the characteristic variety

Proposition

Let $\mathcal{B} \subseteq \mathcal{A}$ be a subarrangement and let π be a neighbourly partition for \mathcal{B} such that dim $(\mathfrak{I}(\pi)) \ge 2$. Then the component passing through $(1, \ldots, 1)$ of the variety in $(\mathbb{C}^*)^{n+1}$ defined by the ideal

$$\mathtt{I}(\pi) + (\mathtt{T}_j - 1 \mid j \notin \mathtt{supp}(\pi))$$

is a non-local component of $\mathcal{V}(\mathcal{A})$. All non-local components of $\mathcal{V}(\mathcal{A})$ passing through $(1, \ldots, 1)$ arise from subarrangements of \mathcal{A} this way.

If $\mathcal{B} = \mathcal{A}$, we call the component *essential*.

Example: A3



- 4 local components
- 1 component with equations
 - $t_1-t_4, \ t_2-t_5, \ t_3t_4t_5-1$

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Example: B3x

This example was discovered by Suciu in 2002



- 7 local components
- 5 components of type A3
- 1 translated component with equations

$$\begin{array}{rl} t_6+1, & t_2-t_3, & t_1-t_4, \\ t_5t_7-1, & t_4t_7+t_3, \\ t_3t_5+t_4, & t_4^2-t_5, \\ t_3t_4+1, & t_3^2-t_7 \end{array}$$



Let \mathcal{A} be an arrangement of n + 1 lines in $\mathbb{P}^2(\mathbb{R})$ and let $\mathcal{M} = \mathcal{M}(\mathcal{A}_{\mathbb{C}}).$

- Alexander matrix from a presentation of $\pi_1(M)$;
- refined Salvetti complex (Salvetti-Settepanella '07; Gaiffi-Salvetti '09): algebraic complex that computes the homology of M with local coefficients.

Theorem

Let $\vartheta_2(t)$ be the 2-boundary map of the refined Salvetti complex that computes $H_*(M;\mathbb{C}_t).$ Then

 $\mathcal{V}(\mathcal{A}_{\mathbb{C}}) = \{ t \in (\mathbb{C}^*)^n \mid \, \mathsf{rk}([\mathfrak{d}_2](t)) < n-1 \}.$



LINE ARRANGEMENTS Computing the characteristic variety

• Compute the primary decomposition of the ideal of all $(n-1) \times (n-1)$ minors.



LINE ARRANGEMENTS

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- Compute one (n−1) × (n−1) minor at a time, factoring it and computing "partial components".



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Remark

LINE ARRANGEMENTS

The number of minors is

$$n \cdot \binom{\nu}{n-1}$$

where
$$\nu = \sum_{P \in Sing(A)} (m(P) - 1)$$
.

CPU total time for computation of $\mathcal{V}(aB3x)$:

Computer	Processor	Benchmark	Time
sedna lab6	Intel Atom N550 AMD A8-3850 APU	235 995	2 h 38 min 46 s 26 min 14 s
lnx1	Intel Xeon E5-2643 v4	2060	13 min 19 s

(benchmark: www.cpubenchmark.net, single thread, last checked on June 11, 2018)





- 11 local components
- 10 components of type A3
- 4 translated components with equations

A new algorithm

We developed a new algorithm that computes the characteristic variety through a series of *bifurcations*.



Comparison of the two algorithms

Computer	Old algorithm	New algorithm
sedna	2 h 38 min 46 s	2 min 21 s
lab6	26 min 14 s	27 S
lnx1	13 min 19 s	14 S

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$\mathcal{R}(12)$



- 16 local components
- 23 comp. of type A3
- 1 comp. of type NonPappus
- ▶ 3 trans. comp. of type B3x
- 1 comp. with equations

$$t_8-t_{11},\ t_7-t_{10},\ t_6-t_9,$$

$$t_5 - t_{11}, t_4 - t_{10}, t_3 - t_9, t_2 - t_9, t_1 - t_{10}, t_9 t_{10} t_{11} - 1$$

▶ 2 trans. comp. with equations

$$\begin{array}{l} t_{10}+t_{11}, \ t_9+t_{11}, \ t_8+t_{11}, \\ t_7+t_{11}, \ t_6+t_{11}, \ t_5-t_{11}, \\ t_4-t_{11}, \ t_3-t_{11}, \ t_2-t_{11}, \\ t_1-t_{11}, \ t_{11}^2-t_{11}+1 \end{array}$$

A(12, 2)



- 14 local components
- ► 35 comp. of type A3
- 2 comp. of type B3
- 1 comp. of type NonPappus
- ▶ 10 trans. comp. of type B3x
- ▶ 8 trans. comp. of type A(10, 2)
- 4 trans. comp. of type $\mathcal{A}(11, 1)$
- > 2 trans. comp. with equations

$$\begin{array}{c} t_{10}-t_{11}, \ t_9+1, \ t_8-t_{11}+1, \\ t_7+1, \ t_6-t_{11}, \ t_5-t_{11}, \\ t_4-t_{11}+1, \ t_3-t_{11}+1, \\ t_2-t_{11}+1, \ t_1-t_{11}+1, \\ t_{11}^2-t_{11}+1 \end{array}$$

Double points partitions

Let \mathcal{A} be a line arrangement in $\mathbb{P}^2(\mathbb{R})$. The double points graph $\Gamma(\mathcal{A})$ is the graph defined as follows:

- its vertex set is $\{H \mid H \in A\}$;
- there is an edge $\{H_1, H_2\}$ iff $H_1 \cap H_2$ is a double point.

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- there is an edge $\{H_1, H_2\}$ iff $H_1 \cap H_2$ is a double point.

Definition

The double points partition of \mathcal{A} is the partition $\Pi_{\mathcal{A}}$ induced by the connected components of $\Gamma(\mathcal{A})$.



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Double points partitions

The definition of $\mathfrak{I}(\pi)$ does not require that π is neighbourly.

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Double points partitions

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Proposition

For all arrangements \mathcal{A} of which we computed $\mathcal{V}(\mathcal{A})$ except one, if $\Pi_{\mathcal{A}}$ is the double points partition of \mathcal{A} , the essential translated components of $\mathcal{V}(\mathcal{A})$ (if exist) appear as the zero locus of one ideal of the primary decomposition of (the radical of) $\mathfrak{I}(\Pi_{\mathcal{A}})$.



Example: DPP of B3x



The primary decomposition of $\mathcal{I}(145|237|6|\infty)$ is $I_1 \cap I_2$, where

$$I_{1} = \begin{pmatrix} t_{6} + 1, t_{2} - t_{3}, t_{1} - t_{4}, t_{5}t_{7} - 1, t_{4}t_{7} + t_{3}, \\ t_{3}t_{5} + t_{4}, t_{4}^{2} - t_{5}, t_{3}t_{4} + 1, t_{3}^{2} - t_{7} \end{pmatrix}$$

$$I_{2} = \begin{pmatrix} t_{6} - 1, t_{2} - t_{3}, t_{1} - t_{4}, t_{5}t_{7} - 1, t_{4}t_{7} - t_{3}, \\ t_{3}t_{5} - t_{4}, t_{4}^{2} - t_{5}, t_{3}t_{4} - 1, t_{3}^{2} - t_{7} \end{pmatrix}$$

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Counterexample: A(11, 1)



- 12 local components
- > 25 comp. of type A3
- 1 comp. of type B3
- 1 comp. of type NonPappus
- ▶ 8 trans. comp. of type B3x
- 4 trans. comp. of type $\mathcal{A}(10,2)$
- 2 trans. comp. with equations

$$\begin{array}{c} t_9-t_{10}, \ t_8+1, \ t_7-t_{10}+1, \\ t_6+1, \ t_5-t_{10}, \ t_4-t_{10}, \\ t_3-t_{10}+1, \ t_2+t_{10}, \\ t_1-t_{10}+1, \ t_{10}^2-t_{10}+1 \end{array}$$

Counterexample: A(11, 1)

$$\Pi_{\mathcal{A}} = (2\,4\,5\,6\,7\,8\,9\,10\,|\,1\,|\,3\,|\,\infty)$$

The two translated components are 0-dimensional, whereas all irreducible components of $\mathcal{Z}(\mathcal{I}(\Pi_{\mathcal{A}}))$ are 1-dimensional. However, the two points *do* belong to $\mathcal{Z}(\mathcal{I}(\Pi_{\mathcal{A}}))$.

Counterexample: A(11, 1)

$$\Pi_{\mathcal{A}} = (245678910|1|3|\infty)$$

The two translated components are 0-dimensional, whereas all irreducible components of $\mathcal{Z}(\mathcal{I}(\Pi_{\mathcal{A}}))$ are 1-dimensional. However, the two points *do* belong to $\mathcal{Z}(\mathcal{I}(\Pi_{\mathcal{A}}))$.

It turns out that the two translated components appear in the primary decomposition of

$$\mathfrak{I}(\Pi_{\mathcal{A}}) + (t_2 t_4 t_9 - 1, t_2 t_5 t_{10} - 1)$$

which is the ideal generated by all the polynomials $\prod t_i - 1$ associated with singular points with multiplicity at least three.

The torus

Definition

A complex algebraic torus ${\mathfrak T}$ is an affine variety isomorphic to $({\mathbb C}^*)^n.$

Definition

A character of \mathcal{T} is a group homomorphism $\chi: \mathcal{T} \to \mathbb{C}^*$ that is a morphism of algebraic varieties. The set of characters of \mathcal{T} is a group $X^*(\mathcal{T})$ isomorphic to \mathbb{Z}^n .

Definition

A one-parameter subgroup of \mathfrak{T} is a group homomorphism $\lambda \colon \mathbb{C}^* \to \mathfrak{T}$ that is a morphism of algebraic varieties. The set of one-parameter subgroups of \mathfrak{T} is a group $X_*(\mathfrak{T})$ isomorphic to \mathbb{Z}^n .

Toric arrangements

Definition

A layer in $\ensuremath{\mathbb{T}}$ is a set of the form

$$\mathcal{K}(\Gamma, \phi) \coloneqq \{ t \in \mathfrak{T} \mid \chi(t) = \phi(\chi) \text{ for all } \chi \in \Gamma \}$$

where $\Gamma < X^*(\mathfrak{T})$ is a split direct summand and $\phi\colon \Gamma \to \mathbb{C}^*$ is a homomorphism.

Definition

A toric arrangement \mathcal{A} in \mathcal{T} is a finite set of layers in \mathcal{T} .



Wonderful models

Wonderful models have been introduced by De Concini and Procesi in 1995 for subspace arrangements.

Definition

A projective wonderful model Y_A for $\mathcal{M}(\mathcal{A})$ is a smooth projective variety containing $\mathcal{M}(\mathcal{A})$ as a dense open set and such that the complement $Y_A \smallsetminus \mathcal{M}(\mathcal{A})$ is a divisor with normal crossings and smooth irreducible components.



TORIC ARRANGEMENTS

Building the wonderful model



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TORIC ARRANGEMENTS

Building the wonderful model



A toric variety can be obtained from a *polyhedral rational fan* Δ in $V \coloneqq X_*(\mathfrak{T}) \otimes \mathbb{R}$.



Equal sign bases

There is a pairing $\langle \cdot, \cdot \rangle$: $X^*(\mathfrak{T}) \times X_*(\mathfrak{T}) \to \mathbb{Z}$.

Definition

Let Δ be a fan in V. A character $\chi \in X^*(\mathfrak{T})$ has the equal sign property with respect to Δ if, for every cone $C \in \Delta$, either $\langle \chi, c \rangle \ge 0$ for all $c \in C$ or $\langle \chi, c \rangle \le 0$ for all $c \in C$.

Definition

Let Δ be a fan in V and let $\mathcal{K}(\Gamma, \varphi)$ be a layer. A basis (χ_1, \ldots, χ_m) for Γ is an equal sign basis with respect to Δ if χ_i has the equal sign property for all $i = 1, \ldots, m$.

 $X_{\mathcal{A}}$ is "good" if it is projective, smooth and every layer of \mathcal{A} has an equal sign basis w.r.t. the fan associated with $X_{\mathcal{A}}$.



Two algorithms

For each $\mathcal{K}_i = \mathcal{K}(\Gamma_i, \phi_i) \in \mathcal{A}$, let $\chi_{i,1}, \dots, \chi_{i,s_i}$ be a \mathbb{Z} -basis of Γ_i and let

$$\Xi = \bigcup_{\mathcal{K}_i \in \mathcal{A}} \{\chi_{i,1}, \dots, \chi_{i,s_i}\}.$$

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Two algorithms

For each $\mathcal{K}_i = \mathcal{K}(\Gamma_i, \phi_i) \in \mathcal{A}$, let $\chi_{i,1}, \dots, \chi_{i,s_i}$ be a \mathbb{Z} -basis of Γ_i and let

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1. Start with a smooth, projective fan and subdivide it so that the final fan is equal sign.



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$$\Xi = \bigcup_{\mathcal{K}_i \in \mathcal{A}} \{ \chi_{i,1}, \dots, \chi_{i,s_i} \}.$$

- 1. Start with a smooth, projective fan and subdivide it so that the final fan is equal sign.
- 2. Start with an equal sign, projective fan and subdivide it so that the final fan is smooth.



DCG algorithm

- 1. Start with the orthant fan (corresponding variety: $(\mathbb{P}^1)^n$).
- 2. Choose a vector $\chi \in \Xi$.
- 3. Repeat until there are no "bad" cones:
 - 3.1 Create the list of bad cones.
 - 3.2 Choose a bad cone and subdivide it.
- 4. Repeat for all vectors in $\boldsymbol{\Xi}$, using the last computed fan as input.

Subdivision can be done in such a way that each computed fan is still smooth and projective.



Smooth algorithm

- 1. Start with the fan generated by the vectors orthogonal to the ones in $\boldsymbol{\Xi}.$
- 2. For each non-smooth cone of the fan, subdivide it in two cones. This can be done in such a way that at least one of them is smooth, and eventually both of them are.



TORIC ARRANGEMENTS

Example



$$\Xi = \{(-1,3), (-1,4)\}$$

Cohomology of the wonderful model

De Concini and Gaiffi (2018): presentation of $H^*(Y_{\mathcal{A}}; \mathbb{Z})$.

- Cohomology ring $H^*(X_{\mathcal{A}}; \mathbb{Z})$
- Well-connected building set G
 - ▶ Poset of layers C(A)
- Adapted bases for the elements of $\mathcal{C}(\mathcal{A})$

 $H^*(Y_{\mathcal{A}};\mathbb{Z})\simeq H^*(X_{\mathcal{A}};\mathbb{Z})[\mathsf{T}_G\mid G\in\mathfrak{G}]\big/\text{some relations}$

Cohomology of the toric variety

Let X be a smooth complete toric variety with associated fan Δ and let \mathcal{R} be the set of the primitive rays of Δ .

$$\mathrm{H}^*(X;\mathbb{Z}) \simeq \mathbb{Z}[C_r \mid r \in \mathfrak{R}] \big/ (\mathrm{I}_{\mathsf{SR}} + \mathrm{I}_{\mathsf{L}})$$

where

•
$$I_{SR}$$
 is the Stanley-Reisner ideal
 $I_{SR} \coloneqq (C_{r_1} \cdots C_{r_k} | r_1, \dots, r_k \text{ do not belong to a cone of } \Delta);$

I_L is the linear equivalence ideal

$$I_{\mathsf{L}} \coloneqq \left(\sum_{\boldsymbol{r} \in \mathcal{R}} \langle \beta, \boldsymbol{r} \rangle C_{\boldsymbol{r}} \; \middle| \; \beta \in X^{*}(\mathfrak{T}) \right).$$

Poset of layers

Definition

Let \mathcal{A} be a toric arrangement in the torus \mathfrak{T} . The poset of layers $\mathfrak{C}(\mathcal{A})$ is the set of all the connected components of the non-empty intersections of the layers of \mathcal{A} , partially ordered by reverse inclusion. It includes \mathfrak{T} as the intersection of zero layers.

To compute $\mathcal{C}(\mathcal{A})$ we use an algorithm by Lenz (2017). However he considers arrangements in the *real compact* torus $(S^1)^n$ instead of the complex algebraic torus $(\mathbb{C}^*)^n$. This is not a problem, as our definition is more general.



Building sets

Let $\mathcal{C}_0(\mathcal{A}) \coloneqq \mathcal{C}(\mathcal{A}) \smallsetminus \{\mathfrak{T}\}$. For the sake of simplicity, assume that all the non-empty intersections of the layers of \mathcal{A} are connected.

Definition

A subset $\mathcal{G} \subseteq \mathcal{C}_0(\mathcal{A})$ is a building set for \mathcal{A} if for each layer $\mathcal{K} \in \mathcal{C}_0(\mathcal{A}) \smallsetminus \mathcal{G}$ the minimal elements of the set $\{G \in \mathcal{G} \mid G \supseteq \mathcal{K}\}$ intersect transversally and their intersection is \mathcal{K} .

Definition

A building set \mathcal{G} is well-connected if for any subset $\{G_1, \ldots, G_k\} \subseteq \mathcal{G}$, if the intersection $G_1 \cap \cdots \cap G_k$ has two or more connected components, then each of them belongs to \mathcal{G} .



Adapted bases

For every pair $(M, G) \in \mathcal{C}(\mathcal{A}) \times \mathcal{C}(\mathcal{A})$ with $G \subseteq M$, we choose a basis $(\beta_1, \ldots, \beta_s)$ for Γ_G such that $(\beta_1, \ldots, \beta_k)$, with $k \leq s$, is a basis for Γ_M . We omit the details here.

From these bases, we compute polynomials $P_G^M \in H^*(X_A; \mathbb{Z})[Z]$, from which the relations for $H^*(Y_A; \mathbb{Z})$ can be obtained.





A two-dimensional example



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TORIC ARRANGEMENTS

A two-dimensional example (cont'd)



 $H^4(Y_{\mathcal{A}};\mathbb{Z})\simeq \mathbb{Z}, \qquad H^2(Y_{\mathcal{A}};\mathbb{Z})\simeq \mathbb{Z}^{14}, \qquad H^0(Y_{\mathcal{A}};\mathbb{Z})\simeq \mathbb{Z}.$

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A five-dimensional example

$$\begin{split} & \mathcal{K}_1 = \{(t_1, \dots, t_5) \in (\mathbb{C}^*)^5 \mid t_1 t_4^{-1} = 1\}, \\ & \mathcal{K}_2 = \{(t_1, \dots, t_5) \in (\mathbb{C}^*)^5 \mid t_2 t_5^{-1} = 1\}, \\ & \mathcal{K}_3 = \{(t_1, \dots, t_5) \in (\mathbb{C}^*)^5 \mid t_3 t_4 t_5 = 1\}. \end{split}$$



$$\begin{split} \mathrm{H}^{10}(\mathrm{Y}_{\mathcal{A}};\mathbb{Z}) &\simeq \mathbb{Z},\\ \mathrm{H}^{8}(\mathrm{Y}_{\mathcal{A}};\mathbb{Z}) &\simeq \mathbb{Z}^{29},\\ \mathrm{H}^{6}(\mathrm{Y}_{\mathcal{A}};\mathbb{Z}) &\simeq \mathbb{Z}^{132},\\ \mathrm{H}^{4}(\mathrm{Y}_{\mathcal{A}};\mathbb{Z}) &\simeq \mathbb{Z}^{132},\\ \mathrm{H}^{2}(\mathrm{Y}_{\mathcal{A}};\mathbb{Z}) &\simeq \mathbb{Z}^{29},\\ \mathrm{H}^{0}(\mathrm{Y}_{\mathcal{A}};\mathbb{Z}) &\simeq \mathbb{Z}. \end{split}$$



Future directions

Line arrangements: investigate the link between the ideals generated by multiple points and the components of the characteristic variety.



Future directions

- Line arrangements: investigate the link between the ideals generated by multiple points and the components of the characteristic variety.
- Toric arrangements: generalize the construction, allowing arbitrary building sets as well as toric arrangements according to De Concini and Gaiffi's definition.





Thank you for your attention.

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