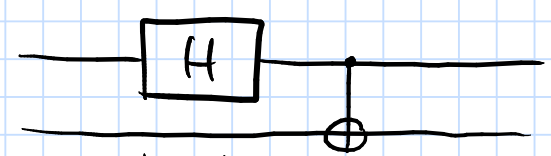


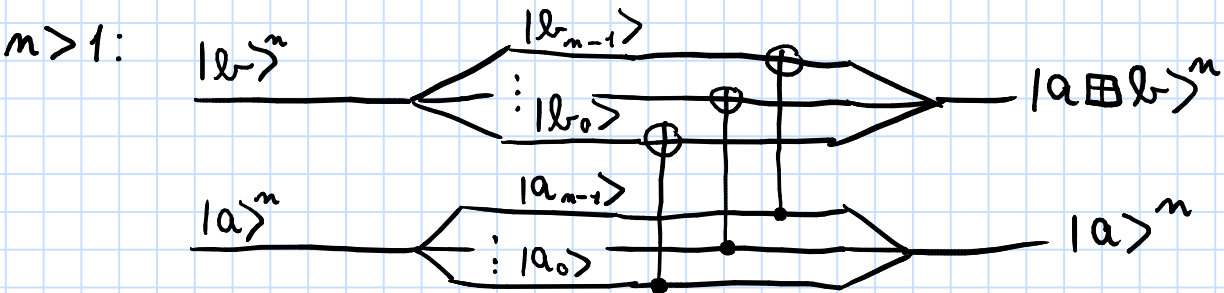
Exc.: consider the circuit  $U = A^\dagger(X)H \otimes I$ :



$U|00\rangle = |\Phi^+\rangle$ , do the other cases.

2)  $U_f: \mathbb{H}^m \otimes \mathbb{H}^m \rightarrow \mathbb{H}^m \otimes \mathbb{H}^m$ ,  $f: \mathbb{N} \rightarrow \mathbb{N}$ .  
 $|a\rangle \otimes |b\rangle \mapsto |a\rangle \otimes |b \oplus f(a)\rangle$

Ex.:  $U_{id} = U_{\oplus}$ .  $m=1$ :  $|a\rangle \otimes |b\rangle \mapsto |a\rangle \otimes |a \oplus b\rangle$



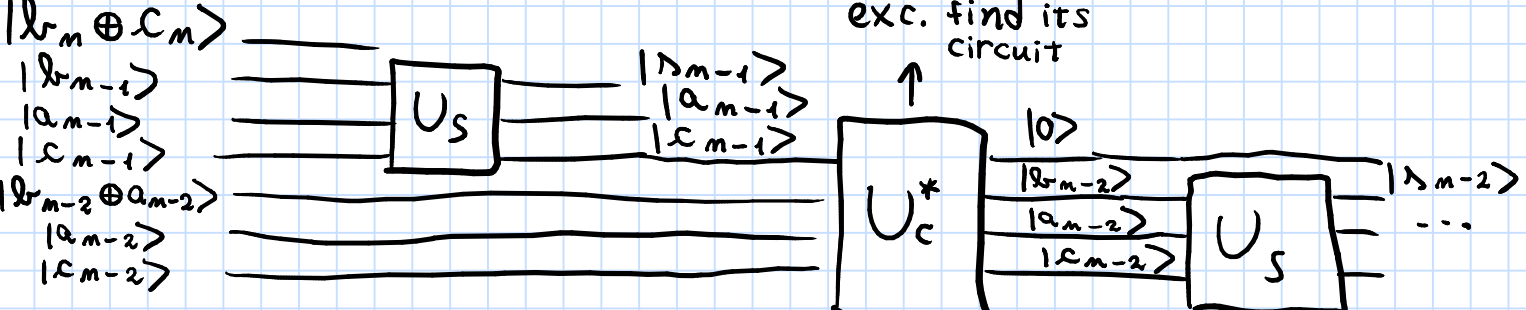
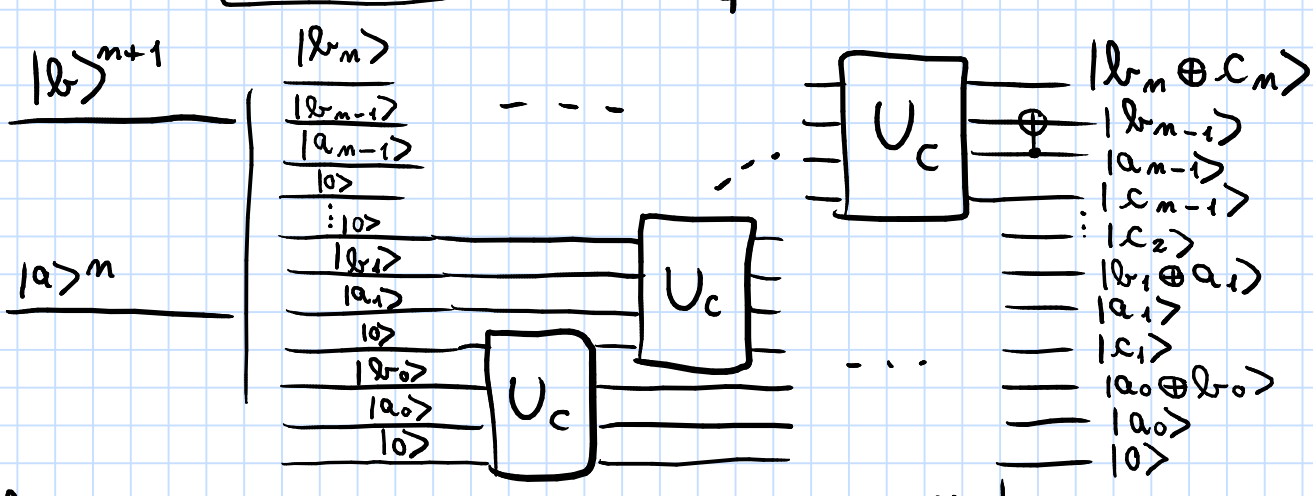
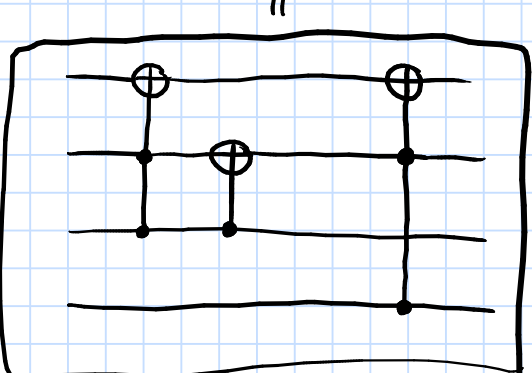
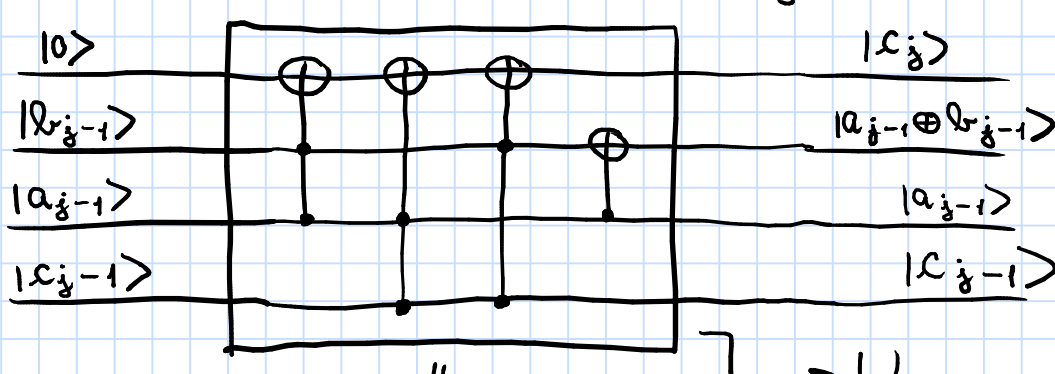
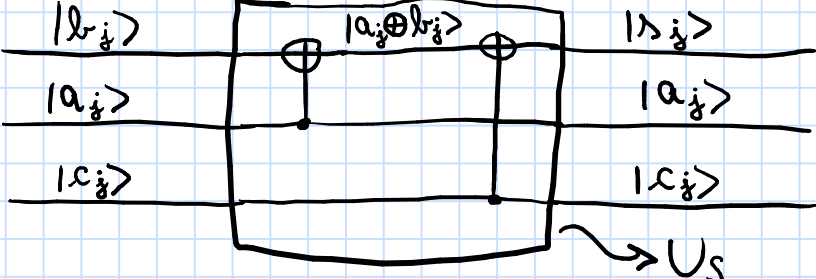
Ex.:  $U_+ : \mathbb{H}^m \otimes \mathbb{H}^{m+1} \rightarrow \mathbb{H}^m \otimes \mathbb{H}^{m+1}$   
 $|a\rangle^m \otimes |b\rangle^{m+1} \mapsto |a\rangle^m \otimes |a+b\rangle^{m+1} \forall a, b \in \{0, 1, \dots, 2^m - 1\}$

$a = \sum_{j=0}^{m-1} 2^j a_j$ ,  $b = \sum_{j=0}^{m-1} 2^j b_j$ ,  $a_j, b_j \in \{0, 1\}$ .

$a+b = \sum_{j=0}^{m-1} 2^j s_j + 2^m c_m$  where we define

$\forall j=0, 1, \dots, m$   $c_j := \begin{cases} 0 & j=0 \\ (a_{j-1} b_{j-1}) \oplus (a_{j-1} c_{j-1}) \oplus (b_{j-1} c_{j-1}) & j=1, \dots, m \end{cases}$

$\forall j=0, 1, \dots, m-1$   $s_j = a_j \oplus b_j \oplus c_j$ .



Exc.: try with  $\tilde{U}_c$

The construction requires  $O(m)$  "elementary gates".

Extensions:  $U_{+ \text{ mod } N} : \mathbb{H}^m \otimes \mathbb{H}^m \rightarrow \mathbb{H}^m \otimes \mathbb{H}^m$   
 $|a\rangle^m \otimes |b\rangle^m \mapsto |a\rangle^m \otimes |a+b \text{ mod } N\rangle^m$

$\forall a, b \in \{0, 1, \dots, N-1\}$  provided that  $2^m \geq N$ .  
 Define  $\mathbb{H}^{<N} \subseteq \mathbb{H}^m$  as  $\text{span}\{|a\rangle^m | a \in \{0, 1, \dots, N-1\}\}$ ;  
 then  $U_{+ \text{ mod } N} : \mathbb{H}^{<N} \otimes \mathbb{H}^{<N} \rightarrow \mathbb{H}^{<N} \otimes \mathbb{H}^{<N}$ .

Similarly,  $U_{- \text{ mod } N} = U_{+ \text{ mod } N}^*$ ,  $U_{c \text{ mod } N} : |a\rangle \otimes |b\rangle \mapsto |a\rangle \otimes |b + a c \text{ mod } N\rangle$ ,

$U_{b \text{ mod } N} : \mathbb{H}^{<N} \rightarrow \mathbb{H}^{<N}$ ,  $|a\rangle \mapsto |b a \text{ mod } N\rangle$ ,  $a, b \in \{0, 1, \dots, N-1\}$ .

Quantum Fourier Transform

Recall: finite F.T.:  $F: (\mathbb{C}_x)_{x=0}^{N-1} \mapsto (\mathbb{C}_\xi)_{\xi=0}^{N-1} = \frac{1}{\sqrt{N}} \sum_{x=0}^{N-1} e^{2\pi i \xi x / N} \cdot C_x$ .

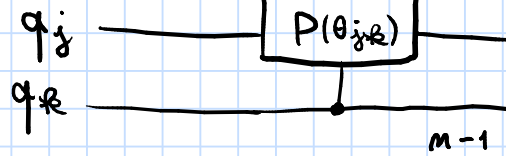
Fix  $N=2^m$ ,  $F: \mathbb{C}^{2^m} = \mathbb{H}^{\otimes m} \rightarrow \mathbb{H}^{\otimes m}$  unitary s.t.  
 $F|x\rangle^m = \frac{1}{\sqrt{2^m}} \sum_{\xi=0}^{2^m-1} e^{2\pi i x \xi / 2^m} |\xi\rangle^m$ ,  $F = (F_{\xi x})_{\xi, x=0, 1, \dots, 2^m-1}$ ,  $F_{\xi x} = \frac{e^{2\pi i \xi x / 2^m}}{\sqrt{2^m}}$ .

$C_{xy} = \begin{cases} 1 & x=y \\ 0 & x \neq y \end{cases}$  How to implement F?

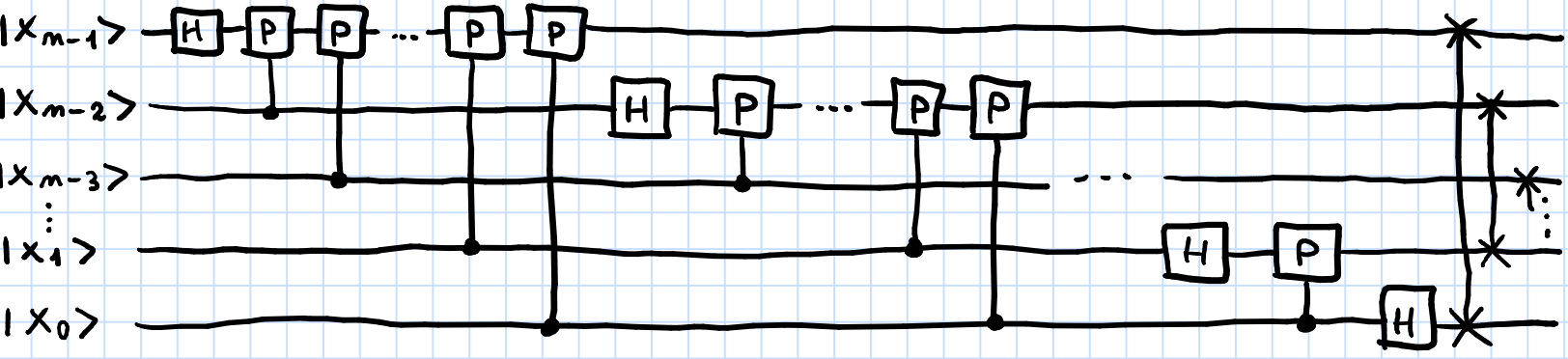
Lemma: we have for  $|x\rangle^m \in \mathbb{H}^m$ ,  $x = \sum_{j=0}^{m-1} x_j 2^j$ ,  
 $F|x\rangle^m = \frac{1}{2^{m/2}} \bigotimes_{j=0}^{m-1} (|0\rangle + e^{2\pi i (0, x_j x_{j-1} \dots x_0)} |1\rangle)$  where  
 $0, x_j x_{j-1} \dots x_0 = \frac{x_j}{2} + \frac{x_{j-1}}{2^2} + \dots + \frac{x_0}{2^{j+1}}$ .

Proof: do the calculations.  $\square$

Define  $P_{jk}$  "controlled phase shift by  $\theta_{jk} = \frac{\pi}{2^{j-k}}$  with control at qubit  $k$ , act on target qubit  $j$ ",  $j > k$ .



Then one has  $F = \overset{\text{swap}}{\sum}^{(m)} \prod_{j=0}^{m-1} \left( \prod_{k=0}^{j-1} P_{jk} \right) H_j$ .



It has  $O(m^2)$  gates.