

Oss.: data  $g(x)$  ol. in  $\sigma_1 < \sigma < \sigma_2$  e continua in  $\sigma_1 \leq \sigma \leq \sigma_2$  e  
 t.c.  $g(\sigma + it) \xrightarrow{|t| \rightarrow \infty} 0 \quad \forall \sigma_1 \leq \sigma \leq \sigma_2$ , si può dimostrare che  
 $f(x) = \frac{1}{2\pi i} \int_{\sigma-i\infty}^{\sigma+i\infty} g(s) x^{-s} ds$  non dipende da  $\sigma$  e si ha  $g(s) = \hat{f}(s)$ .

Teorema (Landau, vedi cap. 3 del Titchmarsh): se si hanno  $\phi(t), \theta(t)$  t.c.  
 $\phi(t) \nearrow +\infty, \theta(t) \searrow, 0 < \theta(t) \leq 1, \frac{\phi(t)}{\theta(t)} = o(e^{\phi(t)})$  e per  $1 - \theta(t) \leq \sigma \leq 2$   
 si ha  $\zeta(s) \ll e^{\phi(t)}$ , allora

$$\beta < 1 - \frac{\theta(2t+1)}{C_0 \phi(2t+1)}$$

In particolare,  $\phi(t) = \log t, \theta(t) = \frac{1}{2} \Rightarrow \beta < 1 - \frac{1}{2C_0 \log(2t+1)}$ .

Littlewood:  $\theta(t) = \frac{(\log \log t)^2}{\log t}, \phi(t) = A \log \log t \Rightarrow \beta < 1 - \frac{c \log \log T}{\log T}$ .

Anche (Vinogradov):  $\theta(t) = \frac{A}{(\log t)^{2/3-2\epsilon}}, \phi(t) = (\log t)^\epsilon,$

$$\zeta(s) \ll \exp((\log t)^\epsilon) \Rightarrow \beta < 1 - \frac{c}{(\log t)^{2/3-\epsilon}}$$

$$1) \sum_{N < m \leq 2N} e^{2\pi i \alpha m} \ll \frac{1}{\|\alpha\|}, \quad 0 < \alpha < 1, \|\alpha\| = \min\{|\alpha - m|, m \in \mathbb{N} \cup \{0\}\}$$

$$2) \sum_{N < m \leq 2N} \Lambda(m) e^{2\pi i \alpha m} \ll \left( \frac{N}{\sqrt{q}} + \sqrt{Nq} + N^{3/4} \right) \log^4 N \text{ dove}$$

$$|\alpha - \frac{a}{q}| \leq \frac{1}{q^2} \quad \hookrightarrow 2N+1 = p_1 + p_2 + p_3 \text{ (1930, Vinogradov)}$$

$$3) \sum_{N < m \leq 2N} m^{-it} \text{ Vinogradov, 1958} \rightsquigarrow \zeta(\sigma + it) \text{ quando } 1 - \sigma \ll \frac{1}{(\log t)^{2/3-\epsilon}}$$

$$4) \sum_{N < m \leq 2N} \frac{\Lambda(m)}{m^{it}} \text{ ? Turán: } \sum_{N < m \leq 2N} \frac{\Lambda(m)}{m^{it}} \ll \frac{N}{t^\epsilon} \text{ per } N \geq t^a \Rightarrow$$

$$\Rightarrow \beta < 1 - \frac{\epsilon}{a}$$

## Caratteri di Dirichlet

$q$  intero  $> 1$ : i caratteri modulo  $q$  sono i caratteri del gruppo  
 abeliano  $(\mathbb{Z}/q\mathbb{Z})^*$ .

$$\chi: (\mathbb{Z}/q\mathbb{Z})^* \rightarrow \mathbb{C}, \quad \phi(q) \text{ caratteri.}$$

$\chi_0$  carattere principale.  
 $(\chi_1 \cdot \chi_2)(m) = \chi_1(m) \cdot \chi_2(m).$   
 $\chi^{-1}(m) = \overline{\chi(m)}.$

$\chi: \mathbb{Z} \rightarrow \mathbb{C}$ :

- 1)  $\chi(mm) = \chi(m)\chi(m) \quad \forall m, m \in \mathbb{Z}$
- 2)  $\chi$  ha periodo  $q$
- 3)  $|\chi(m)| = 1$  se  $(m, q) = 1$ , o altrimenti.

Caso  $p^a, p > 2$  primo.  $(\mathbb{Z}/p^a\mathbb{Z})^*$  è ciclico,  $g$  generatore.

$\omega$  t.c.  $\omega^{\phi(p^a)} = 1$ .  $m$  intero,  $\nu(m)$  t.c.  $g^{\nu(m)} = [m]_{\mathbb{Z}/p^a\mathbb{Z}}$

$$\chi_p(m) = \omega^{\nu(m)} = e^{\frac{2\pi i m \nu(m)}{\phi(p^a)}}$$

Il caso  $(\mathbb{Z}/2\mathbb{Z})^*$  è banale.

Caso  $2^2$ :  $(\mathbb{Z}/4\mathbb{Z})^*$ ,  $\omega = 1 \quad \chi_0$   
 $\omega = -1 = e^{\frac{2\pi i}{2}}$

$$\chi_4(m) = \begin{cases} 1 & m \equiv 1 \pmod{4} \\ -1 & m \equiv 3 \pmod{4} \\ 0 & m \text{ pari} \end{cases}$$

$g$  generatore,  $g^2 \equiv [m]$

Caso  $2^a, a \geq 3$ :  $(\mathbb{Z}/2^a\mathbb{Z})^* \cong \mathbb{Z}/2\mathbb{Z} \times \mathbb{Z}/2^{a-2}\mathbb{Z}$ .

$$m \equiv (-1)^{\nu_2(m)} 5^{\nu_1(m)} \pmod{2^a}, \quad \chi(m) = e^{\pi i m \nu_0 + \frac{2\pi i m \nu_1}{2^{a-2}}}$$

$m \leftrightarrow$  carattere,  $\nu \leftrightarrow m$ .

Caso  $q = 2^a p_1^{a_1} \dots p_n^{a_n}$

$$\chi_q(m) = e^{\pi i m \nu_0 + \frac{2\pi i m \nu_1}{2^{a-2}} + \frac{2\pi i m \nu_2}{\phi(p_1^{a_1})} + \dots + \frac{2\pi i m \nu_n}{\phi(p_n^{a_n})}}$$

$$4) \sum_{m=1}^q \chi(m) = \begin{cases} 0 & \text{se } \chi \neq \chi_0 \\ \phi(q) & \text{se } \chi = \chi_0 \end{cases}$$

$$5) \sum_{\chi \text{ mod } q} \chi(m) = \begin{cases} 0 & \text{se } m \not\equiv 1 \pmod{q} \\ \phi(q) & \text{se } m \equiv 1 \pmod{q} \end{cases}$$

$$\overline{\chi}(a) = \chi^{-1}(a) = \chi(a^{-1}). \quad 5) \Rightarrow \text{se } (a, q) = 1,$$

$$\frac{1}{\phi(q)} \sum_{\chi \text{ mod } q} \overline{\chi}(a) \chi(m) = \begin{cases} 1 & \text{se } m \equiv a \pmod{q} \\ 0 & \text{altrimenti} \end{cases}$$

Def.: sia  $\chi$  carattere mod  $q, \chi \neq \chi_0$ . Si dice funzione  $L$  relativa  
 al carattere  $\chi$  la somma della seguente serie:

$$L(s, \chi) = \sum_{m=1}^{\infty} \frac{\chi(m)}{m^s} \quad (\sigma > 1).$$

Prop.: se  $\chi \neq \chi_0$  la serie converge uniformemente per  $\sigma \geq \epsilon > 0$ .

$$\text{Dim.: } \sum_{m \leq N} \frac{\chi(m)}{m^s} = \left( \sum_{m \leq N} \chi(m) \right) N^{-s} + \int_1^N \left( \sum_{m \leq u} \chi(m) \right) \frac{du}{u^{s+1}}.$$

$$\left| \sum_{m \leq N} \chi(m) \right| \leq q, \quad \sigma > 1 \Rightarrow$$

$$\Rightarrow L(s, \chi) = \int_1^{+\infty} \left( \sum_{m \leq u} \chi(m) \right) \frac{du}{u^{s+1}}, \text{ e questo converge unif.}$$

in  $\sigma \geq \epsilon > 0$ .  $\square$

Oss.:  $\chi = \chi_0, \sum_{m=1}^N \frac{\chi_0(m)}{m^s}$ ; si ha  $\sum_{m \leq N} \chi_0(m) \leq \left( \left\lfloor \frac{N}{q} \right\rfloor + 1 \right) \phi(q) \Rightarrow$

$\Rightarrow$  non si ottiene conv. unif. per  $\sigma \geq \epsilon > 0$ .

Oss.: per l'identità di Eulero si ha

$$L(s, \chi) = \prod_p \left( 1 - \frac{\chi(p)}{p^s} \right)^{-1} \text{ per } \sigma > 1 \quad \forall \chi \text{ mod } q.$$

$$L(s, \chi_0) = \prod_p \left( 1 - \frac{\chi_0(p)}{p^s} \right)^{-1} = \prod_{p \nmid q} \left( 1 - \frac{1}{p^s} \right)^{-1} =$$

$$= \prod_p \left( 1 - \frac{1}{p^s} \right)^{-1} \prod_{p|q} \left( 1 - \frac{1}{p^s} \right) = \prod_{p|q} \left( 1 - \frac{1}{p^s} \right) \zeta(s) \Rightarrow$$

$\Rightarrow L(s, \chi_0)$  è mero. con polo semplice in  $s=1$  con  
 residuo  $\prod_{p|q} \left( 1 - \frac{1}{p} \right) = \frac{\phi(q)}{q}$ .

Altri zeri:  $p^{-s} = 1$ , cioè  $\sigma=0$  e  $t = \frac{2k\pi}{\log p}$ .

Oss.:  $\log L(s, \chi) = - \sum_p \log \left( 1 - \frac{\chi(p)}{p^s} \right) = \sum_p \sum_{m=1}^{\infty} \frac{\chi(p^m)}{m p^{ms}}, \quad \sigma > 1$

$$\frac{L'}{L}(s, \chi) = - \sum_{m=1}^{\infty} \frac{\chi(m) \Lambda(m)}{m^s}$$

$$\frac{L'}{L}(s, \chi_0) = - \sum_{\substack{m=1 \\ (m, q)=1}}^{\infty} \frac{\chi_0(m) \Lambda(m)}{m^s} = - \sum_{m=1}^{\infty} \frac{\Lambda(m)}{m^s} - \sum_{p|q} \frac{\log p}{p^{s+1}}$$

$$\log L(s, \chi) = \sum_p \frac{\chi(p)}{p^s} + O\left( \sum_p \sum_{m \geq 2} \frac{1}{p^{sm}} \right) = \sum_p \frac{\chi(p)}{p^s} + O\left( \sum_p \frac{1}{p^{\sigma(p^2-1)}} \right) =$$

$$= \sum_p \frac{\chi(p)}{p^s} + O(1).$$

$$\sum_{\chi \text{ mod } q} \log L(s, \chi) = \sum_{\chi \text{ mod } q} \sum_p \frac{\chi(p)}{p^s} + O(\phi(q)). \text{ Prendendo } s \text{ reale,}$$

$$\log \prod_{\chi \text{ mod } q} |L(s, \chi)| = \sum_p \left( \frac{1}{p^s} + \sum_{\chi \neq \chi_0} \frac{\chi(p)}{p^s} \right) + O(\phi(q))$$